

Algorithmische Graphentheorie

SS 09

The backtracking approach of Brown

Procedure: Color(k)

{First possibility: assign v_k sequentially to one of the already existing color classes, if feasible.}

$NrColors[k] := NrColors[k - 1];$

{ determine the colors feasible for v_k : }

for $i = 1$ to $NrColors[k]$ **do**

$FeasColor[k, i] := TRUE;$

end for

for $i: v_i \in \Gamma(v_k)$ **do**

if $i < k$ **then**

$FeasCol[k, c(i)] := FALSE;$

end if

end for

$i = 1;$

while $(i \leq NrColors[k]) \wedge (NrColors[k] < f)$ **do**

if $FeasColor[k, i]$ **then**

$c[k] := i;$

if $k < n$ **then**

 Color(k+1)

else

 {found a better coloring!}

$f := NrColors[k];$

end if

end if

$i := i + 1;$

end while

Procedure Color(k) - continued

{Second possibility: start up a new color class for v_k .}

```
if  $NrColors[k] + 1 < f$  then  
     $NrColors[k] := NrColors[k] + 1;$   
     $c[k] := NrColors[k];$   
    if  $k < n$  then  
        Color(k+1)  
    else  
        {found a better coloring!}  
         $f := NrColors[k];$   
    end if  
end if
```

main program

```
 $f := n;$       { Initialisation of upper bound on  $\chi(G)$  }  
order the nodes in a reasonable ordering  $v_1, v_2, \dots, v_n;$   
 $c[1] := 1;$   
 $NrColors[1] := 1;$   
Color(2);  
output  $f;$ 
```